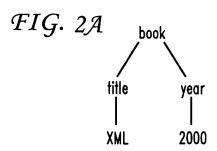
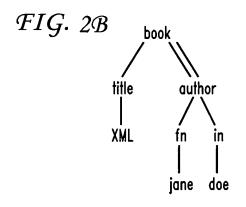
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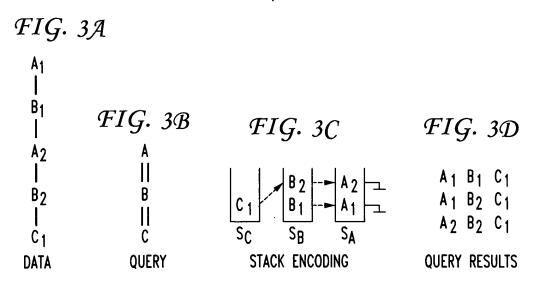
1/17 FIG. 1 book (1,1:50,1) allauthors chapter title year (1,64:93,2) (1,61:63,2) (1,2:4,2) (1,5:60,2)(1,62,3)(1,6:20,3) (1,3,3) XML author author 2000 author title section (1,65:67,3) (1,68:78,3) (1,7:9,4)(1,69:71,4)XML fn In In · fn In head (1,66,4)doe john doe jane **Origins** jane poe (1,8,5) (1,11,5) (1,26,5) (1,43,5) (1,46,5) (1,70,5)





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2/17



```
Algorithm PathStack(q)
01 while \neg \operatorname{end}(q)
         qmin = getMinSource(q)
02
03
         for q_i in subtreeNodes(q) // clean stacks
            while (\neg empty(Sqi) \land topB(Sqi) < nextL(Tqmin))
04
05
               pop(Sq_i)
         moveStreamToStack(Tq_{\min}, Sq_{\min}, pointer to
06
                                           top(S_{parent}(q_{min})))
        if (isleaf(q_{min}))
07
            showSolutions(Sqmin,1)
08
09
            pop (S_{q_{\min}})
Function end(q)
    return \forall q_i \in \text{subtreeNodes}(q) : \text{isLeaf}(q_i) \Rightarrow \text{eof}(Tq_i)
Function getMinSource(q)
    return q_i \in \text{subtreeNodes}(q) such that \text{nextL}(Tq_i)
       is minimal
Procedure moveStreamToStack(Tq,Sq,p)
01 \operatorname{push}(S_q,(\operatorname{next}(T_q),p))
02 advance(T_{a})
                                                 PathStack
```

3/17 FIG. 4A START NO STREAM Q END NOT EMPTY YES GET NODE WITH MINIMUM 102 ORDERING VALUE QI TO PROCESS REMOVE PARTIAL ANSWERS - 104 FROM STACK THAT CANNOT EXTEND TO FULL ANSWERS AUGMENT PARTIAL ANSWERS WITH Qi **- 106** 108 IS NO Qi A LEAF NODE YES PRODUCE SOLUTIONS **- 110**

4/17

```
Procedure showSolutions(SN,SP)
// Assume, for simplicity, that the stack of the query
// nodes from the root to the current leaf node we
// are interested in can be accessed as S[1],...,S[n].
// Also assume that we have a global array index[1..n]
      of pointers to the stack elements.
// index[i] represents the position in the i'th stack that
// we are interested in for the current solution, where
// the bottom of each stack has position 1
// Mark we are interested in position SP of stack SN.
01 \text{ index}[SN] = SP
02 if (SN == 1) // we are in the root
     // output solutions from the stacks
      output (S[n].index[n],...,S[1].index[1])
04
    else // recursive call
05
       for i = 1 to S[SN].index[SN].pointer_to_parent
06
         showSolutions(SN - 1,i)
07
                          Procedure showSolutions
```

FIG. 6

	CASE 1	CASE 1	CASE 1	CASE 1
PROPERTY	X.R <y.l< td=""><td>X.L<y.l X.R>Y.R</y.l </td><td>X.L>Y.L X.R<y.r< td=""><td>X.L>Y.R</td></y.r<></td></y.l<>	X.L <y.l X.R>Y.R</y.l 	X.L>Y.L X.R <y.r< td=""><td>X.L>Y.R</td></y.r<>	X.L>Y.R
SEGMENTS	X Y	Y Y	X Y	<u>X</u>
TREE	ROOT Y	ROOT	ROOT	ROOT
CASES FOR PathStack AND TwigStack				

METHOD AND SYSTEM FOR PATTERN MATCHING HAVING HOLISTIC TWIG JOINS Nicolas Bruno et al. Application No. 10/747,847

5/17

```
Algorithm PathMPMJ(q)
01 while (\neg \operatorname{eof}(T_q) \land (\operatorname{isRoot}(q)))
                                 \mathsf{nextL}(q) \ < \mathsf{nextR}(\mathsf{parent}(q))))
     for (q_i \in \text{subtreeNodes}(q)) // advance descendants
02
         while (\text{nextL}(q_i) < \text{nextL}(\text{parent})q_i)))
03
             advance(T_{q_i})
04
         PushMark(T_{q_i})
05
    if (isLeaf(q)) // solution in the streams' heads outputSolution()
06
      else PathMPMJ(child(q))
07
08 advance(Tq) 09 for (q_i \in \text{subtreeNodes}(q)) // backtrack descendants
         PopMark(T_{q_i})
10
                                                        PathMPMJ
```

6/17

```
Algorithm TwigStack(q)
    // Phase 1
    while \neg end(q)
        q_{act} = getNext(q)
02
        if (\neg isRoot(q_{act}))
03
           cleanStack(parent(act), next(q_{act}))
04
        if (isRoot(q_{act}) V \neg empty(S_{parent}(q_{act})))
05
06
           cleanStack(q_{act}, next(q_{act}))
           moveStreamToStack(Tq_{act}, pointer to
07
                                    top(Sparent(q_{act})))
           if (isLeaf(q_{act}))
80
            showSolutionWithBlocking(Sqact, 1)
09
            pop(Sq_{act})
10
        else advance(Tq_{act})
11
       // Phase 2
12 mergeAllPathSolutions()
Function getNext(q)
     if (isLeaf(q) return q)
01
02 for q_i in children(q)
      n_i = getNext(q_i)
03
      if (n_i \neq q_i) return n_i
     n_{min} = minarg_{n_i} nextL(T_{n_i})
     n_{max} = maxrarg_{n_i} nextL(T_{n_i})
     while (nextR(Tq) < nextL(Tn_{max}))
07
         advance(Tq)
80
09
     if (\text{nextL}(Tq) < \text{nextL}(T_{n_{min}})) return q
     else return nmin
Procedure cleanStack(S, actL)
     while (\neg empty(S) \land (topR(S) < actL))
         pop(S)
02
                                                  TwigStack
```

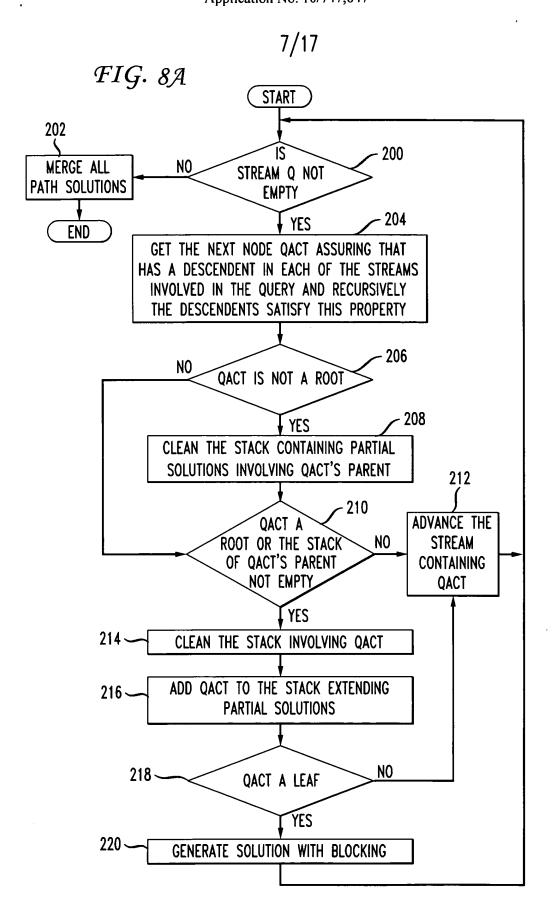
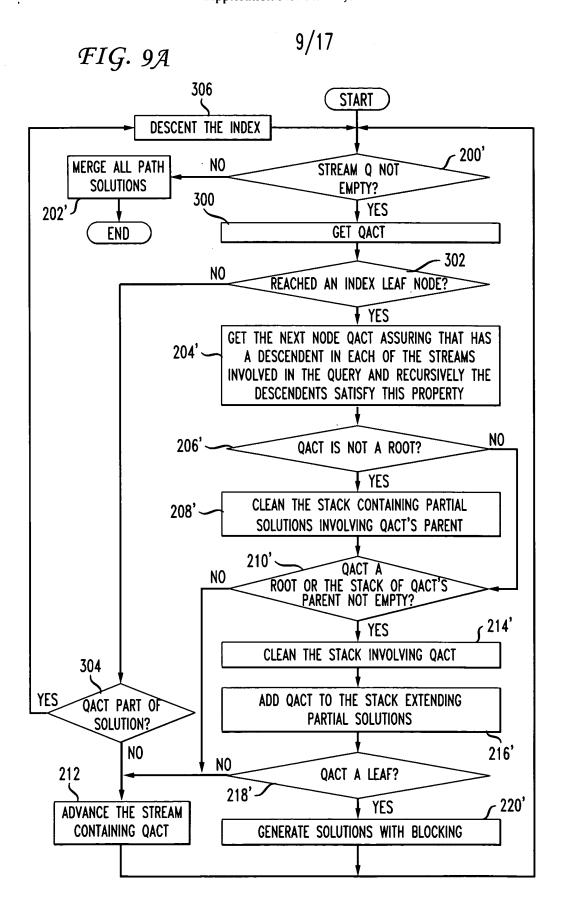


FIG. 9

8/17

```
Algorithm TwigStackXB(q)
01 while \neg \operatorname{end}(q)
02
        q_{act} = getNext(q)
        if (isPlainValue(Tq_{act}))
(03)
           if (\neg isRoot(q_{act}))
04
              cleanStack(parent (q_{act}), next(q_{act}))
05
           if (isRoot(q_{act}) V \neg empty(Sparent(<math>q_{act})))
06
              cleanStack(q_{act}, next(q_{act}))
07
              moveStreamToStack(Tq_{act}, pointer to
08
                                      top(S_{parent(q_{act})}))
              if (isLeaf(q_{act}))
09
               showSolutionsWithBlocking(Sq_{act},1)
10
11
               pop(Sq_{act})
           else advance(Tq_{act})
12
        else if (\neg isRoot(q_{act}) \land empty(Sparent(q_{act})) \land
(13)
                   nextL(T_{parent}(q_{act})) > nextR(T_{qact}))
         advance(T_{qact}) // Not part of a solution
(14)
         else // Might have a child in some solution
(15)
         drillDown(T_{q_{qct}})
(16)
       // Phase 2
17 mergeAllPathSolutions()
Function getNext(q)
01 if (isLeaf(q) return q)
02 for q_i in children(q)
        n_i = getNext(q_i)
03
 (04) if (q_i \neq n_i \ V \neg isPlainValue(T_{n_i})) return n_i
05 n_{min} = minarg_{n_i} nextL(T_{n_i})
06 n_{max} = maxrarg_{n_i} nextL(T_{n_i})
     while (nextR(Tq) < nextL(Tn_{max}))
08
         advance(Tq)
     if (nextL(Tq) < nextL(T_{nmin})) return q
09
    else return nmin
Procedure cleanStack(S, actL)
     while (\neg empty(S) \land (topR(S) < actL))
02
          pop(S)
                                                   TwigStack
```



10/17

FIG. 10

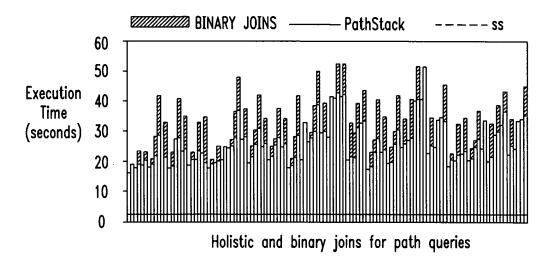
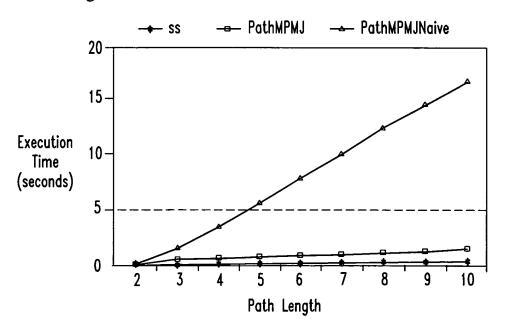


FIG. 11



11/17

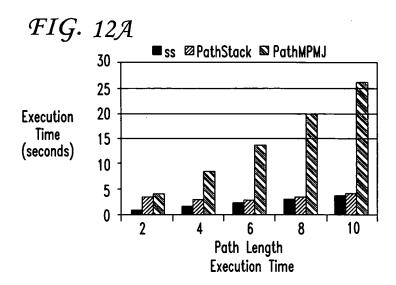
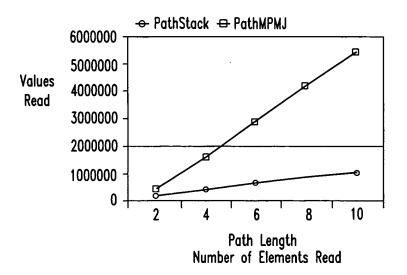


FIG. 12B



12/17

FIG. 13A

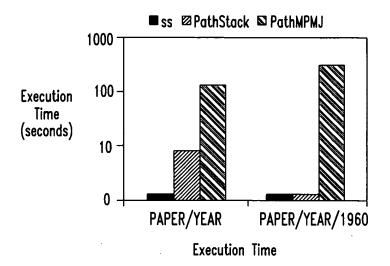
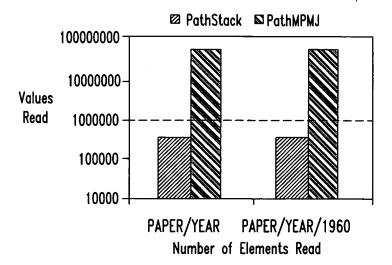
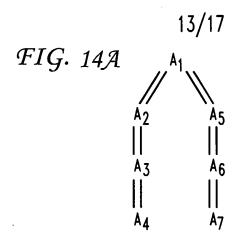
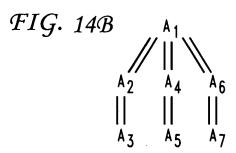
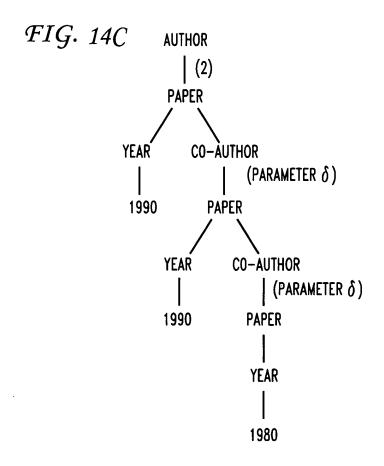


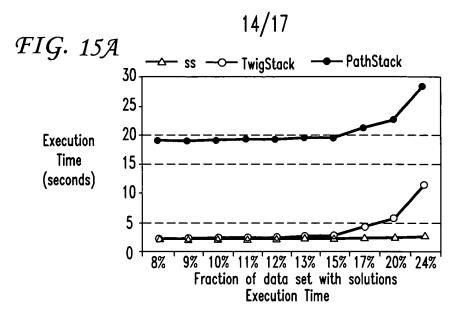
FIG. 13B

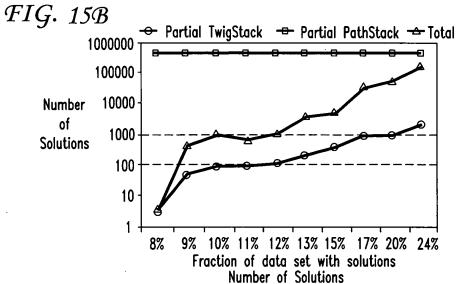


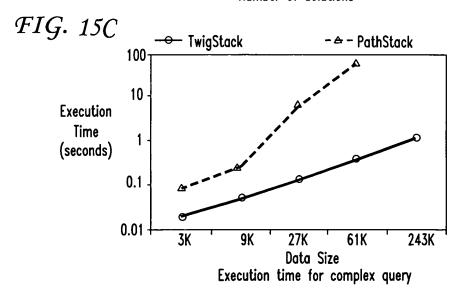


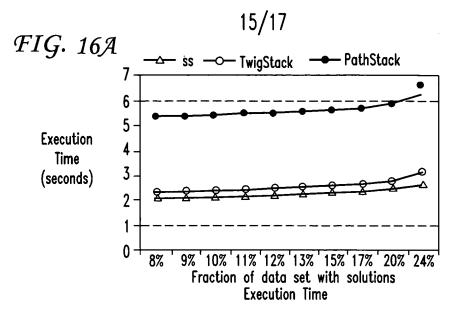


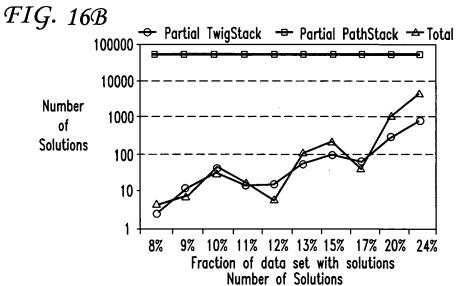


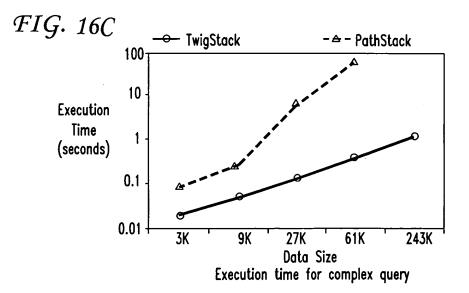












16/17

FIG. 17A

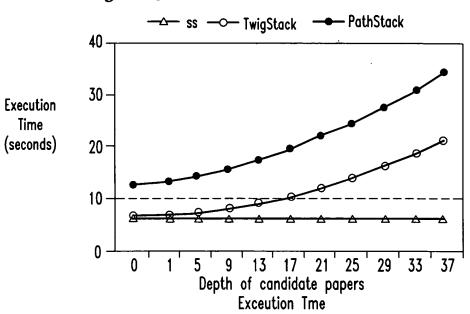


FIG. 17B

